

$$\sum_{j \leq x} \Lambda(j) E(x/j) = \sum_{k \leq x} \log k - 2 \sum_{k \leq x/2} \log k \quad (x \geq 2).$$

Proof. With $E(x) := \sum_{n \leq x} (\mathbf{1} * \nu)(n)$ as above,

$$\begin{aligned} S := \sum_{j \leq x} \Lambda(j) E(x/j) &= \sum_{j \leq x} \sum_{k \leq x/j} \Lambda(j) (\mathbf{1} * \nu)(k) = \sum_{jk \leq x} \Lambda(j) (\mathbf{1} * \nu)(k) \\ &= \sum_{n \leq x} [\Lambda * (\mathbf{1} * \nu)](n) \quad (\text{definition of } *) \\ &= \sum_{n \leq x} (\ell * \nu)(n) \quad (\Lambda * \mathbf{1} = \ell) \\ &= \sum_{j \leq x} \nu(j) \sum_{k \leq x/j} \log k \quad (\text{as above}) \\ &= \sum_{k \leq x} \log k - 2 \sum_{k \leq x/2} \log k \quad (x \geq 2). \quad // \end{aligned}$$

LEMMA 2.

$$\psi(2n) \geq \log \binom{2n}{n}.$$

Proof. Take $x = 2n$ in L. 1. As $E(\cdot) \leq 1$, $S \leq \sum_{j \leq 2n} \Lambda(j) = \psi(2n)$. But

$$\begin{aligned} S &:= \sum_{k \leq x} \log k - 2 \sum_{k \leq x/2} \log k = \sum_{k=n+1}^{2n} \log k - \sum_{k=1}^n \log k = \log \left(\frac{(n+1)(n+2) \dots (2n)}{1.2 \dots n} \right) \\ &= \log \binom{2n}{n}. \quad // \end{aligned}$$

Th. 3 (Chebyshev's Lower Estimates). For $\epsilon > 0$ and x large,

(i) $\psi(x) \geq (\log 2 - \epsilon)x$; (ii) $\theta(x) \geq (\log 2 - \epsilon)x$; (iii) $\pi(x) \geq (\log 2 - \epsilon)li(x)$.

Proof. (i) Let $N := \binom{2n}{n}$ as above. This is the largest of the $2n+1$ terms in the binomial expansion of $(1+1)^{2n} = 2^{2n}$ (by Pascal's triangle), so $2^{2n} \leq (2n+1)N$. So $2n \log 2 \leq \log N + \log(2n+1)$, and Lemma 2 gives

$$\psi(2n) \geq \log N \geq 2n \log 2 - \log(2n+1).$$

Given x , take n with $2n \leq x < 2n + 2$. Then (i) follows as

$$\psi(x) \geq \psi(2n) \geq (x - 2) \log 2 - \log(x + 1).$$

(ii): from (i) as $(\psi(x) - \theta(x))/x \rightarrow 0$;

(iii): from (ii) by the first Theorem of this section. //

Cor. 5. $\pi(x) \geq (\log 2 - \epsilon)x / \log x$.

Proof. $\psi(x) \leq \pi(x) \log x$ (first Prop. of this section and (i). //

THEOREM (Chebyshev, 1849-51) (Mastery Question, 2013).

$$\ell := \liminf \pi(x)/li(x) \leq 1 \leq \limsup \pi(x)/li(x) =: L.$$

In particular, if the limit exists, it is 1 (as in PNT).

Proof. For all $\epsilon > 0$ there exists x_0 such that for $x \geq x_0$

$$\ell - \epsilon \leq \frac{\pi(x)}{x / \log x}, \quad \frac{\pi(x)}{x / \log x} \leq L + \epsilon.$$

For the lower bound, integration by parts gives, as $0 < \pi(u) \leq u$,

$$\sum_{p \leq x} \frac{1}{p} \geq \sum_{x_0 < p \leq x} \frac{1}{p} = \int_{(x_0, x]} \frac{d\psi(u)}{u} = \frac{\pi(x)}{x} - \frac{\pi(x_0)}{x_0} + \int_{x_0}^x \frac{\pi(t)}{t^2} dt,$$

$$\geq -1 + \int_{x_0}^x \frac{\pi(t)}{t^2} dt \geq -1 + (\ell - \epsilon) \int_{x_0}^x \frac{dt}{t \log t} \geq (\ell - \epsilon) \log \log x + O_\epsilon(1)$$

($\int^x dt/(t \log t) = \int^x d \log t / \log t = \log \log x$). But by Mertens' Second Th.,

$$\sum_{p \leq x} 1/p = \log \log x + c_1 + O(1/\log x).$$

Combining,

$$\log \log x + c_1 + O(1/\log x) \geq (\ell - \epsilon) \log \log x + O_\epsilon(1) : \quad 1 \geq \ell - \epsilon.$$

This holds for all $\epsilon > 0$. So $\ell \leq 1$. The upper bound is similar but simpler. //

In 1851, Chebyshev also proved *Bertrand's postulate* of 1845: for any $n \geq 2$ there is a prime p between n and $2n$; see 2013 Problems and Solutions 8 for Erdős' elementary proof of 1932.